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D E C I S I O N
of 15 December 2000

Case Number: T 0189/99 - 3.5.1

Application Number: 93303466.2

Publication Number: 0569218

IPC: G06F 12/02

Language of the proceedings: EN

Title of invention:
Circuit for rotating a digital image

Applicant:
XEROX CORPORATION

Opponent:
-

Headword:
Circuit for rotating a digital image/XEROX CORPORATION

Relevant legal provisions:
EPC Art. 56, 52(1)

Keyword:
"Inventive step (yes)"

Decisions cited:
-

Catchword:
-



Case Number: T 0189/99 - 3.5.1

D E C I S I O N
of the Technical Board of Appeal 3.5.1
of 15 December 2000

Appellant: XEROX CORPORATION
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Decision under appeal: Decision of the Examining Division of the
European Patent Office posted 17 November 1998
refusing European patent application
No. 93 303 466.2 pursuant to Article 97(1) EPC.

Composition of the Board:

Chairman: P. K. J. van den Berg
Members: R. Randes
P. H. Mühlens

Summary of Facts and Submissions

I. This is an appeal against the decision of the Examining Division to refuse European Patent Application 93 303 466.2 because the subject matter of claim 1 lacked inventive step having regard to the prior art acknowledged in figure 6 of the application itself, derived from

D2: EP-A-0 433 645,

and the disclosure of the following document:

D1: EP-A-0 276 110.

II. The Appellant (Applicant) has requested grant of a patent and made an auxiliary request for oral proceedings.

III. The Appellant requests grant of a patent based on the following documents:

Claim:

1 received on 10 October 2000

Description:

Pages 1-3 received on 29 September 1998

Pages 4-5 received on 17 October 1997

Drawings:

Page 2/5 as originally filed

Pages 1/5 and 3/5 to 5/5 received on 17 October 1997.

IV. Claim 1 reads as follows:

"A circuit for addressing one bit at a time in a random access memory device (14) having x address lines divided into column address lines and row address lines, the circuit comprising:

an x bit counter (10) having y least significant bits, z most significant bits, and n intermediate bits, where $y + z + n = x$, and x, y, z and n are positive integers,

a first multiplexer (12) connected to the counter (10) and having A and B inputs and an output, the A and B inputs being connected to the counter (10) with the A input being coupled to the y least significant bits and the B input being coupled to the z most significant bits;

a second multiplexer (13) connected to the counter (10) and having A' and B' inputs and an output, the A' and B' inputs being connected to the counter (10) with the A' input being coupled to the z most significant bits and the B' input being coupled to the y least significant bits; and

a register (11) having $y + z + n$ input and output lines wherein the y least significant of the input lines (A0 - A8) thereof are coupled to the output of the first multiplexer (12), the n intermediate input lines (A9 - A12) thereof are coupled to the n intermediate bits (C9 - C12) of the counter (10) and the z most significant of the input lines (A13 - A21) thereof are coupled to the output of the second multiplexer (13), wherein the multiplexers (12, 13) in a first state couple the A and A' inputs to the register (11) and in a second state couple the B and B', inputs to the register (11);

characterised in that the y least significant bit (A0 - A8) output lines of the register (11) are divided

into lower least significant bit (A0 - A5) output lines and upper least significant bit (A6 - A8) output lines, the z most significant bit (A13 - A21) output lines of the register (11) are divided into lower most significant bit (A13 - A17) output lines and upper most significant bit (A18 - A21) output lines,

said lower least significant bit (A0 - A5) output lines and said lower most significant bit (A13 - A17) output lines being coupled to the column address lines of the random access memory device (14),

said upper least significant bit (A6 - A8) output lines, the intermediate bit (A9 - A12) output lines and said upper most significant bit (A18 - A21) output lines of the register (11) being coupled to the row address lines of the random access memory device (14)."

Reasons for the Decision

1. *Admissibility*

The appeal is admissible.

2. *Background*

Memories for storing and rotating digital images are usually organised as two-dimensional arrays, each memory location being specified by a row address and a column address. If such memories are addressed in such a way that consecutive locations differ only in their column address then faster reading and writing is achievable in what is known as the "fast page mode" of accessing. If an image is merely stored in such a memory and then read out again, then fast page mode addressing is possible for both writing and reading. If

however the image is also to be rotated then data is written into the memory in one "direction", but read out of the memory in another "direction". For instance, when writing data into the memory consecutive locations usually differ only in their column address so that fast page mode accessing is possible. However when reading data out of the memory consecutive locations usually differ only in their row address so that fast page mode accessing is not possible. As a consequence fast page mode accessing is only possible for reading or writing, but not both. This leads to the undesirable result that memory access for image storage occurs at a different rate to memory access for image storage with rotation. The invention overcomes this problem by providing a circuit for addressing which rearranges the memory row and column address lines to change the way in which an image is mapped onto the two-dimensional memory area with the result that fast page mode accessing is, at least partly, possible in both directions.

3. *Amendments*

In addition to editorial amendments and the correction of errors pointed out by the Examining Division in its decision, claim 1 has been amended with respect to the refused version by the addition in the first line of the expression "one bit at a time in" and the addition in the final paragraph of the expression "the intermediate bit (A9 - A12) output lines", these features having been discussed in paragraph 15.2 of the impugned decision. Since these amendments are supported by column 1, lines 50-53 and figure 7 of the published application, the new claim satisfies Article 123(2) EPC. The claim is also clear and supported by the

description.

4. *Novelty*

It is common ground between the Appellant and the Board that the prior art acknowledged in figure 6 of the application represents the closest prior art on file. The Appellant has stated that, even though D2 does not contain a figure corresponding to figure 6 of the application, figure 6 gives a view of the prior art known from D2.

Figure 6 of the application discloses the features set out in the preamble of claim 1 for the case $x=22$, $y=9$, $n=4$, $z=9$. Since the characterising features of claim 1 are not disclosed by figure 6 of the application, the subject matter of claim 1 is novel over this prior art.

D1 concerns the storage of an image having three dimensions (X,Y and S) in a memory which, in response to a row address (RA0-RA7), outputs a complete row of data consisting of several data words. D1 seeks to solve the problem that although data can be read quickly in the X direction by accessing a row at a time, access in the S and Y direction is slow since each bit is in a different row. The solution lies in each row containing data words relating to all three dimensions. Hence data can be read for all directions by reading a row and selecting the appropriate word. This selection is made based on the settings of the B0-B1 and C0-C3 address lines which can thus be regarded as column address lines (see fig. 17). The row and column addresses are generated from external row and column addresses (Y0-Y8 and X0-X8) by the address scramble circuit (50) which is shown in figures 17 and

19 and described by truth tables 6-1 and 6-2.

Hence the address scramble unit (50) of D1 in the X-access mode rearranges memory address lines in a similar way to that set out in the characterising part of claim 1. However D1 does not relate to addressing one memory bit at a time, as required by claim 1. Instead D1 is concerned with addressing one word at a time. Consequently D1 does not mention any intermediate bits and the four least significant external column address bits (X0-X3) are not used by the address scramble unit in X-access mode; see page 13, lines 33-36. Since a circuit for addressing one bit at a time in a random access memory device is not disclosed by D1, the subject matter of claim 1 is also novel over D1.

The Board notes that the remaining document cited in the European Search Report (GB-A-2 092 785) merely reflects background art.

5. *Inventive step*

The subject matter of claim 1 differs from the prior art shown in figure 6 of the application in the features set out in the characterising part of claim 1. These differences have the effect of allowing fast read mode to be used when reading out from the memory both horizontally and vertically and result from solving the technical problem of increasing access speeds in the direction in which fast page mode cannot be used.

The impugned decision states in paragraph 13 that the skilled person confronted with this problem would certainly consult D1 which is in the same technical

field of memories for image processing. Whilst it is true that D1 and D2 belong to the same technical field, the Board is unconvinced for two reasons that the skilled person would regard D1 as offering a solution to the above problem. Firstly, D2, like the circuit set out in claim 1, concerns the addressing of a memory one bit at a time, whilst D1 relates to a memory in which 16-bit words are addressed (see fig. 17; DQ0 - DQ15). Secondly, D1 does not mention intermediate bits and so cannot indicate what should be done with the intermediate bits shown in figure 6 of the application. Claim 1 however makes clear how these intermediate bits are used.

Moreover, as argued by the Appellant, the hypothetical application of the teaching of D1 to the circuit shown in figure 6 of the application does not result in the subject matter of the claim, since the address scramble unit known from D1 converts eighteen external address lines into fourteen internal address lines, thus not using the four least significant X address lines. The result would consequently be the use of only fourteen of the eighteen counter (10) bits to address the memory (14). This possibility does not however fall within claim 1 which requires that the number of memory address lines (x) be equal to the number of bits of the counter (10). The Board can furthermore see no obvious reason for the skilled person to further modify the circuit to also use the four least significant counter bits.

6. The Board concludes that the subject matter of claim 1 involves an inventive step, Articles 52(1) and 56 EPC.

7. *Deficiencies*

Although claim 1 is allowable, the description requires adaptation to the wording of claim 1 to comply with Rule 27(1)(c) EPC. Moreover the reference symbols "A" and "B" used in the claim cannot be found in the figures, particularly figure 7, contrary to Rule 32(2)(i) EPC. The Board consequently remits the case to the Examining Division to attend to these deficiencies.

8. *Oral proceedings*

Since the Appellant's request is allowable, oral proceedings, mentioned at point II above, need not be held.

Order

For these reasons it is decided that:

1. The decision is set aside.
2. The case is remitted to the Examining Division with the order to remove the deficiencies mentioned in paragraph 7 above and to grant a patent on the basis of the Appellant's request (see paragraph III above).

The Registrar:

The Chairman:

M. Kiehl

P. K. J. van den Berg